

# PL Rules for Tableaux(Truth Tree)

## 1 Notation

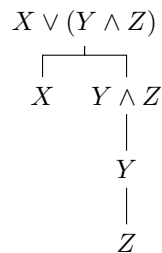
### 1.1 Substitution Instance

Robotnese: for any expression  $\Phi$ , a substitution instance  $\Phi^\beta/\alpha$  refers to the formula obtained by replacing all occurrence of  $\alpha$  in  $\Phi$  by some  $\beta$  not already in  $\Phi$ .

Human-speak: Consider  $\forall xPx$ . Let  $\Phi = Px$ . A substitution instance of  $Px$  could be  $Pc$ , where we  $\alpha = x$  and  $\beta = c$  - All occurrences of  $x$  are replaced with  $c$ .

### 1.2 Root

The lowest node. For instance,  $X \vee (Y \wedge Z)$  is the root for the tree below:



### 1.3 Path

A (directed) path is a non-branching series of connected and nodes. Take the tree above, there two directed paths in this tree. One from the root to  $X$ , and the other one goes from the root to  $Z$ . In our context, we only talk about directed path.

## 2 Tree Rules for PL

$\forall$  **Decomposition** If a non-decomposed (not marked with an X) PL sentence  $\forall\alpha\Psi$  occurs on an open path(i.e., branch), you may:

1. If some constant  $\beta_0$  is already present on that path, add a node on that path with the PL sentence with the substitution instance  $\Phi^{\beta_0}/\alpha$ .
2. For all constants  $\beta_1, ..\beta_n$  on that path, add a new substitution instance as you did for step 1.
3. If no constant exists on that path, pick some  $\beta$  and add a node on that path with the substitution instance  $\Phi^\beta/\alpha$ .
4. In either case, *don't* check the  $\forall\alpha\Phi$ .
5. If a new constant  $\beta_{n+1}$  is later introduced, instantiate  $\forall\alpha\Phi$  by writing  $\Phi^{\beta_{n+1}}/\alpha$ . This might produce infinite trees.

Example:

$$\begin{array}{c}
 \forall xPx(\text{given}) \\
 | \\
 Fa(\text{given}) \\
 | \\
 Pa
 \end{array}$$

**$\exists$  Decomposition** If a non-decomposed (not marked with an X) PL sentence  $\exists\alpha\Psi$  occurs on an open path(i.e., branch), you may:

1. You must pick some  $\beta$  constant that has not occurred *at all* and add a node on that path with the substitution instance  $\Phi^\beta/\alpha$ .
2. Add  $\Phi^\beta/\alpha$  too all open paths.
3. *Do* check the  $\exists\alpha\Psi$ .

Example:

$$\begin{array}{c}
 \exists xPx(\text{given}) \quad X \\
 | \\
 Fa \vee (Gb \vee Mc)(\text{given}) \quad X \\
 \begin{array}{cc}
 | & | \\
 Fa & (Gb \vee Mc) \quad X \\
 | & \begin{array}{cc}
 | & | \\
 Gb & Mc \\
 | & | \\
 Pd & Pd
 \end{array}
 \end{array}
 \end{array}$$

**$\neg\forall$  Decomposition** If a non-decomposed (not marked with an X) PL sentence  $\neg\forall\alpha\Psi$  occurs on an open path(i.e., branch), you may:

1. Add  $\exists\neg\alpha\Psi$  too all paths that contain  $\neg\forall\alpha\Psi$
2. Check the original node

**$\neg\exists$  Decomposition** If a non-decomposed (not marked with an X) PL sentence  $\neg\exists\alpha\Psi$  occurs on an open path(i.e., branch), you may:

1. Add  $\forall\neg\alpha\Psi$  too all paths that contain  $\neg\exists\alpha\Psi$
2. Check the original node